

# On the Sample Complexity of Adversarial Multi-Source PAC Learning

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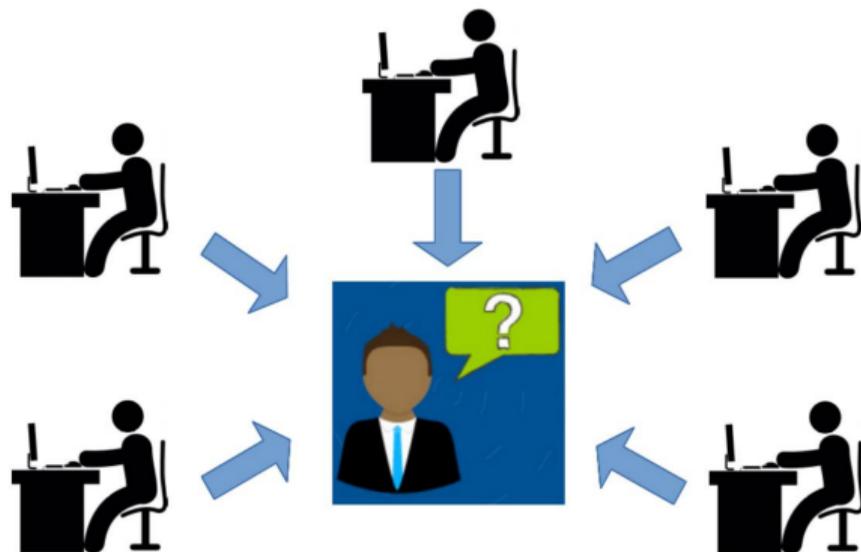
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# Summary

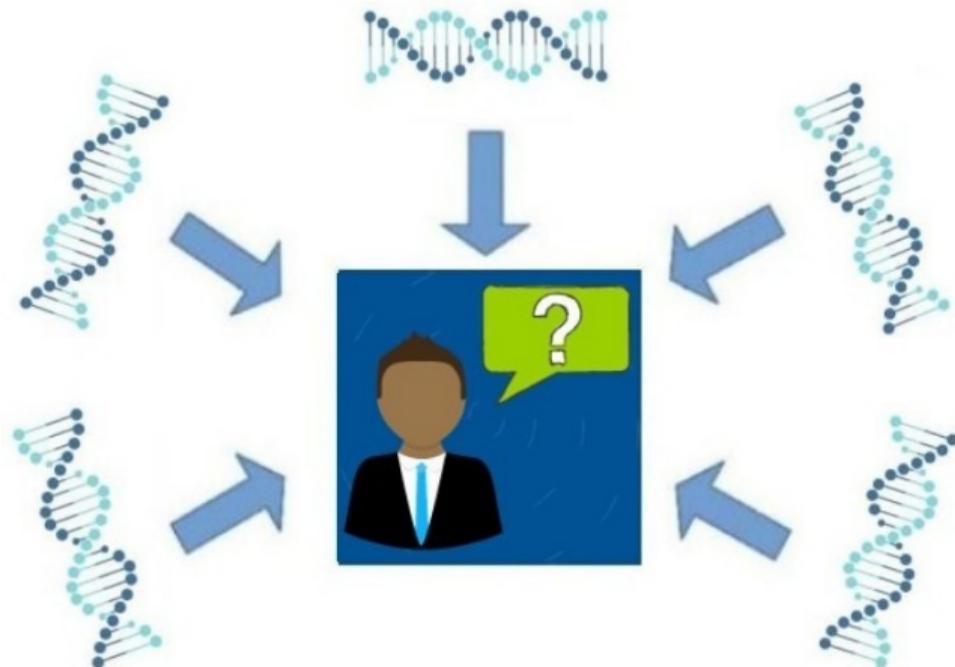
# Learning from untrusted sources

## Crowdsourcing



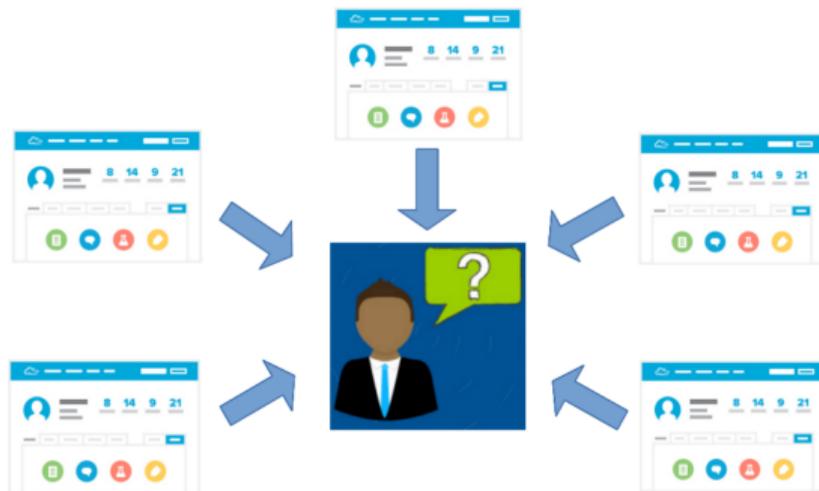
# Learning from untrusted sources

## Using data from multiple labs



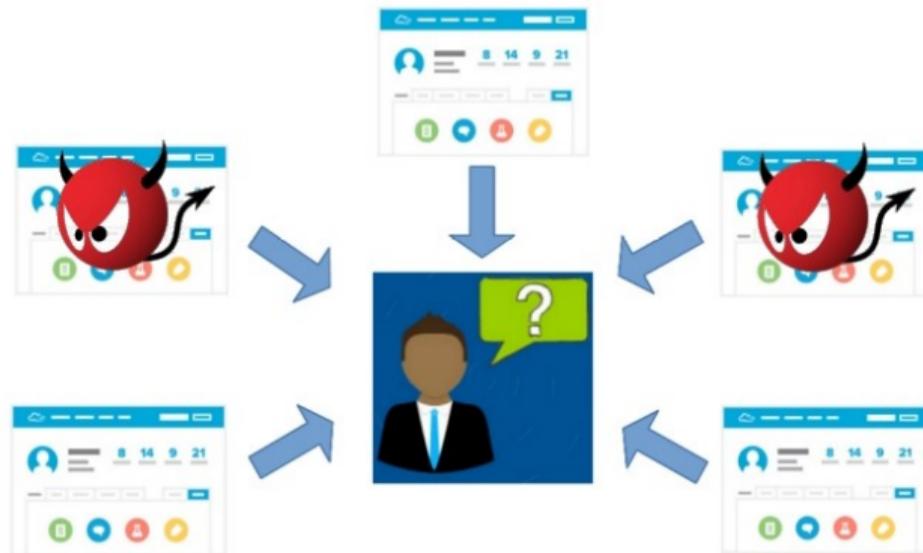
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## Collecting data from online sources



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**How much can be learnt even if some data is corrupted or manipulated?**

# Main contributions

- Rigorous adversarial models and statistical PAC-learnability framework
- Positive results:
  - PAC-learnability is fulfilled (under minimal assumptions)
  - Explicit learning algorithm and rates
- Hardness results:
  - Sample complexity lower bound
  - The learner needs the group structure to achieve PAC-learnability

# Details

# Setup

## Supervised learning scenario

- Input-output space  $\mathcal{X} \times \mathcal{Y}$ , unknown data distribution  $\mathcal{D} \in \mathcal{P}(\mathcal{X} \times \mathcal{Y})$
- Hypothesis space  $\mathcal{H}$ , loss function  $\ell : \mathcal{Y} \times \mathcal{Y} \rightarrow \mathbb{R}^+$
- Want to find  $h \in \mathcal{H}$ , such that  $\mathcal{R}(h) = \mathbb{E}_{\mathcal{D}}(\ell(h(x), y))$  is small

## Learning from multiple sources

- Given: a set of  $N$  datasets  $S = (S_1, \dots, S_N)$
- $m$  labeled points in each:  $S_i = \{(x_{i,j}, y_{i,j})\}_{j=1}^m \stackrel{\text{iid}}{\sim} \mathcal{D}$

# Adversarial model

## Informal description

- An adversary controls an  $\alpha$ -fraction of the sources,  $\alpha < 1/2$
- The adversary can choose the new points with full knowledge of the setup
- The learner does not know which sources are manipulated

## Formal definitions

- $(\mathcal{X} \times \mathcal{Y})^{N \times m}$  - set of all unordered sequences of  $N$  sets of  $m$  points
- A fixed-set adversary is *any function*  $\mathcal{A} : (\mathcal{X} \times \mathcal{Y})^{N \times m} \rightarrow (\mathcal{X} \times \mathcal{Y})^{N \times m}$ , such that:

$$(S'_1, \dots, S'_N) = \mathcal{A}(S_1, \dots, S_N) \text{ satisfies } S'_i = S_i,$$

$\forall i \in C$ , where  $C$  is the set of “clean” sources and  $|C| = (1 - \alpha)N$

# Adversarial PAC-learnability

- A multi-source learner is a function  $\mathcal{L} : (\mathcal{X} \times \mathcal{Y})^{N \times m} \rightarrow \mathcal{H}$
- Focus on fixed  $N$  and  $\alpha$ , while  $m \rightarrow \infty$
- $\mathcal{H}$  is  $\alpha$ -adversarially PAC-learnable if  $\exists m : (0, 1]^2 \rightarrow \mathbb{N}$ , such that for any  $\epsilon, \delta \in (0, 1]$ , whenever  $m \geq m(\epsilon, \delta)$ , with probability at least  $1 - \delta$ :

$$\mathcal{R}(\mathcal{L}(\mathcal{A}(S))) \leq \min_{h \in \mathcal{H}} \mathcal{R}(h) + \epsilon,$$

against *any (fixed-set) adversary of power  $\alpha$*

## Related work

### Learning discrete distributions from untrusted batches

- Unsupervised version of the problem studied in (Qiao and Valiant 2018; Jain et al. 2020)

### Robust PAC learning from a single dataset

- One point per source recovers the malicious noise model (Kearns et al. 1993)
- PAC-learnability is known to be impossible: minimum possible error is  $\alpha/(1 - \alpha)$

### Byzantine-robust distributed optimization

- Practical and robust gradient optimization methods (Yin et al. 2018; Alistarh et al. 2018)
- Convergence analysis under convexity/smoothness assumptions

### Collaborative learning

- Multiple parties learn *one model each*
- Adversarial PAC-learnability provably possible (Blum et al. 2017; Qiao 2018)

# Adversarial PAC-learnability

**Main assumption:**  $\mathcal{H}$  is uniformly convergent

- Given  $m$  samples  $S = \{(x_1, y_1), \dots, (x_m, y_m)\} \stackrel{iid}{\sim} \mathcal{D}$ , with probability at least  $1 - \delta$  over the data :

$$\sup_{h \in \mathcal{H}} |\mathcal{R}(h) - \widehat{\mathcal{R}}(h)| \leq s_{\mathcal{H}, \ell}(m, \delta, S),$$

- $s_{\mathcal{H}, \ell}(m, \delta, S_m) \rightarrow 0$  as  $m \rightarrow \infty$ , for any sequence  $\{S_m\}_{m \in \mathbb{N}}$  with  $S_m \in (\mathcal{X} \times \mathcal{Y})^m$

## Theorem

$\mathcal{H}$  - uniformly convergent  $\implies \mathcal{H}$  - adversarially PAC-learnable.

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## Theorem

$\mathcal{H}$  - uniformly convergent  $\implies \mathcal{H}$  - adversarially PAC-learnable.

Holds even against a stronger adversary that can choose which sources to corrupt

## Sample complexity upper bound

- In many situations  $s_{\mathcal{H},\ell}(m, \delta, S) = \mathcal{O}(1/\sqrt{m})$
- There exists a learning algorithm, such that with probability at least  $1 - \delta$ :

$$\mathcal{R}(\mathcal{L}(\mathcal{A}(S))) - \min_{h \in \mathcal{H}} \mathcal{R}(h) \leq \tilde{\mathcal{O}}\left(\frac{1}{\sqrt{(1-\alpha)Nm}} + \alpha \frac{1}{\sqrt{m}}\right),$$

against any fixed-set adversary<sup>1</sup>

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<sup>1</sup> $\tilde{\mathcal{O}}$  hides constants and logarithmic factors

## Hardness results (for formal statements see paper)

### Sample complexity lower bound

- No learning algorithm can achieve against any adversary error less than:

$$\mathcal{O}\left(\frac{1}{\sqrt{(1-\alpha)Nm}} + \alpha\frac{1}{m}\right)$$

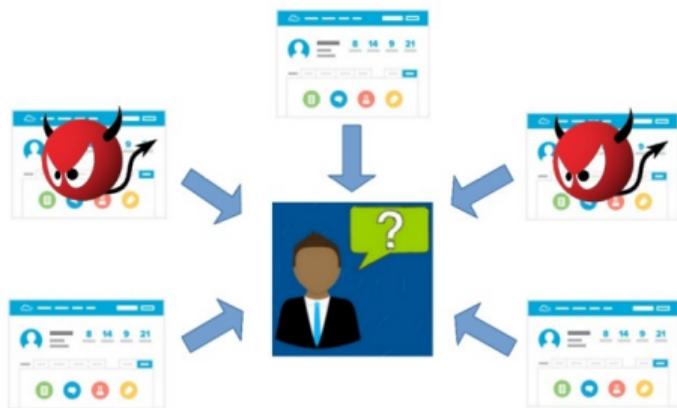
- If  $m$  is constant and  $\alpha > 0$ ,  $N \rightarrow \infty$  does not guarantee PAC-learnability

### The learner has to use the group structure

- No learning algorithm that ignores the group structure can guarantee error less than  $\mathcal{O}(\alpha/(1-\alpha))$

## Summary

- Learning from multiple unreliable sources now commonplace
- Setup modeled as a PAC-learning problem with an adversary
- Group structure enables PAC-learnability, even against a strong adversary
- Describe fundamental limitations on the learner



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Thank you for your attention!

Meet us at the poster session for more details.

## References I

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