

Oracle Efficient Private Non-Convex Optimization

Seth Neel

University of Pennsylvania

Giuseppe Vietri

University of Minnesota

Aaron Roth

University of Pennsylvania

Steven Wu

University of Minnesota

First Setting.

- Consider following general optimization problem defined in terms of dataset $D \in X^n$.

$$\min_{w \in \mathcal{W}} L(D, w)$$

- \mathcal{W} is a **discrete** and bounded set. For this talk assume that $\|\mathcal{W}\|_\infty = 1$
- The loss function $L : X^n \times \mathcal{W} \rightarrow [0,1]$ is **bounded**.

Can we solve solve this problem with differential privacy?

First Contribution

- We propose an Objective Perturbation Algorithm that:
 - Satisfies (ϵ, δ) -differential privacy.
 - W.h.p finds an answer $\hat{w} \in \mathcal{W}$ with error bounded by

$$\left| L(D, \hat{w}) - L(D, w^*) \right| \leq \frac{14\sqrt{2(d+1)\ln(2\beta)}\sqrt{\ln(1/\delta)}}{n\tau\epsilon}.$$

The Normalization Trick

- Let $\|w\|_2 \leq D$, For all $w \in \mathcal{W}$. The normalization function is:

$$\pi(w) = \left(w_1, \dots, w_d, D\sqrt{1 - \|w\|_2/D} \right) \frac{1}{D}$$

- Key idea: Sample random vector η , and Augment objective:

$$\min_{w \in \mathcal{W}} L(D, w) - \langle \eta, \pi(w) \rangle$$

- Solve using non-private optimization oracle.

Stability of the Objective

- Let $w, w' \in \mathcal{W}$. Then

$$\|\pi(w) - \pi(w')\|_2^2 \geq \frac{1}{D^2} \|w - w'\|_2^2$$

- The normalization trick provides stability.

Objective Perturbation.

Algorithm 1: ObjPertDiscrete

Input: Projection function π , $\mathcal{D} = \{l_i\}$, optimization oracle

$n \leftarrow |\mathcal{D}|$;

$\sigma \leftarrow \frac{7LD^2\sqrt{\ln 1/\delta}}{\tau\epsilon}$;

Draw i.i.d random vector $\eta \sim \mathcal{N}(0, \sigma^2)^{d+1}$;

$$w \in \arg \min_{w \in \mathcal{W}_\tau} \left(\frac{L(\mathcal{D}, w) - \langle \eta, \pi(w) \rangle}{n} \right)$$

Output: w

Second Setting

- Consider following general optimization problem defined in terms of dataset $D \in X^n$.

$$\min_{w \in \mathcal{W}} L(D, w)$$

- \mathcal{W} is a **general** decision space. For this talk assume that $\|\mathcal{W}\|_\infty = 1$
- The loss function $L : X^n \times \mathcal{W} \rightarrow [0,1]$ is **G-Lipschitz, convex and bounded.**

Can we solve solve this problem with differential privacy?

Second Contribution

- We propose an Objective Perturbation Algorithm that:
 - Satisfies (ϵ, δ) -differential privacy.
 - W.h.p finds an answer $\hat{w} \in \mathcal{W}$ with error bounded by

$$\left| L(D, \hat{w}) - L(D, w^*) \right| \leq \frac{d^{5/4} G \sqrt{D_2 \log(1/\beta)}}{\sqrt{n\epsilon}}$$

Standard Differential Privacy (DP)

[Dwork et al., 2006]

Two datasets



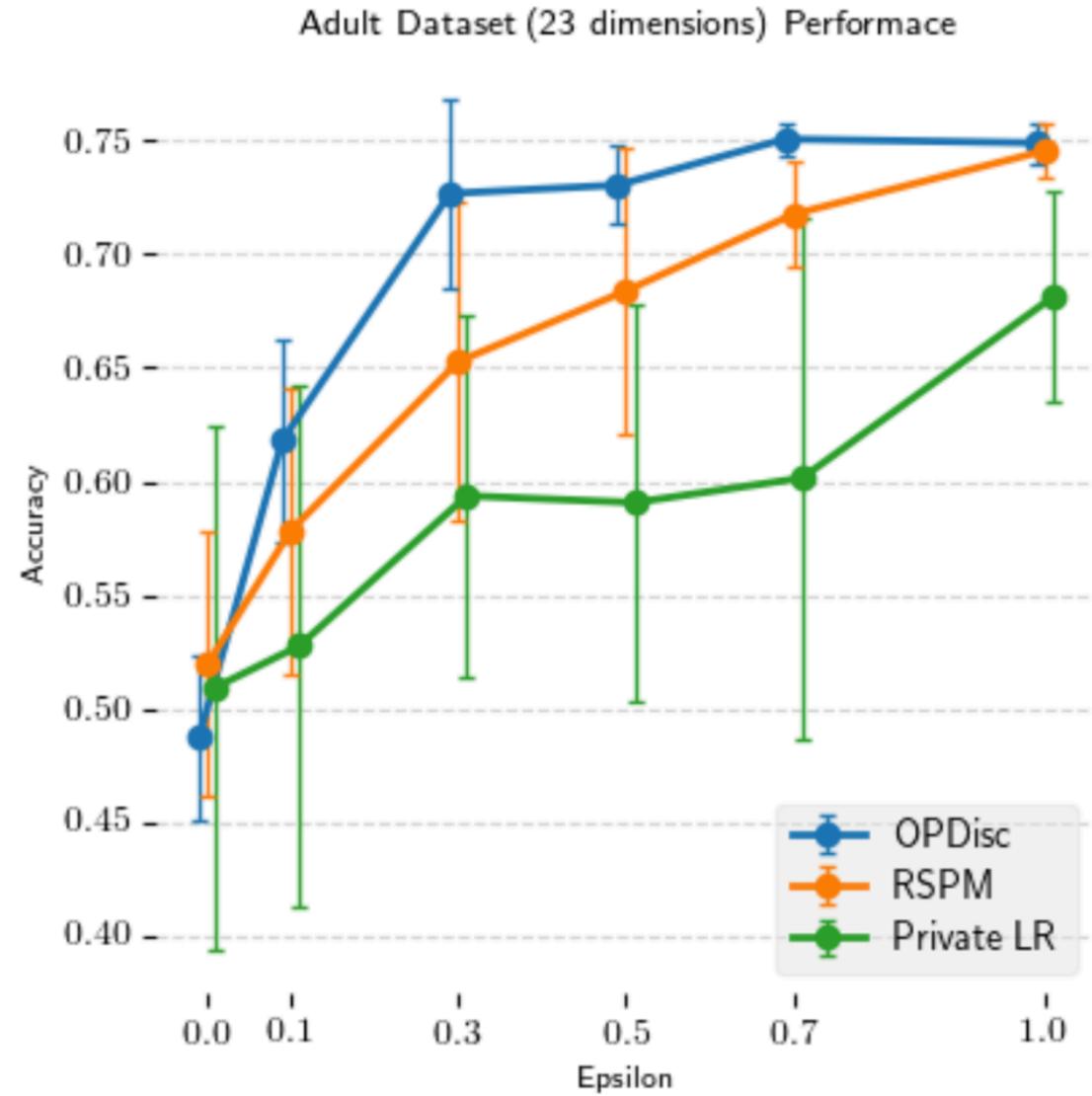
are **neighbors** if they are different on only one row.

Definition: Mechanism M satisfies ϵ -differential privacy if, for all neighboring datasets and for all $r \in \text{range}(M)$

$$\Pr[M(\text{red database}) = r] \leq e^\epsilon \Pr[M(\text{blue database}) = r]$$

Experiments.

Experiment



Conclusion.

- We provide two new private algorithm,